

From JSCert to JSExplain and Beyond

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From JSCert

What is JSCert?

Two JavaScript semantics in Coq

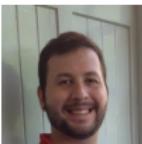
descriptive given a program and a result, say if they are related

executable given a program, compute the result

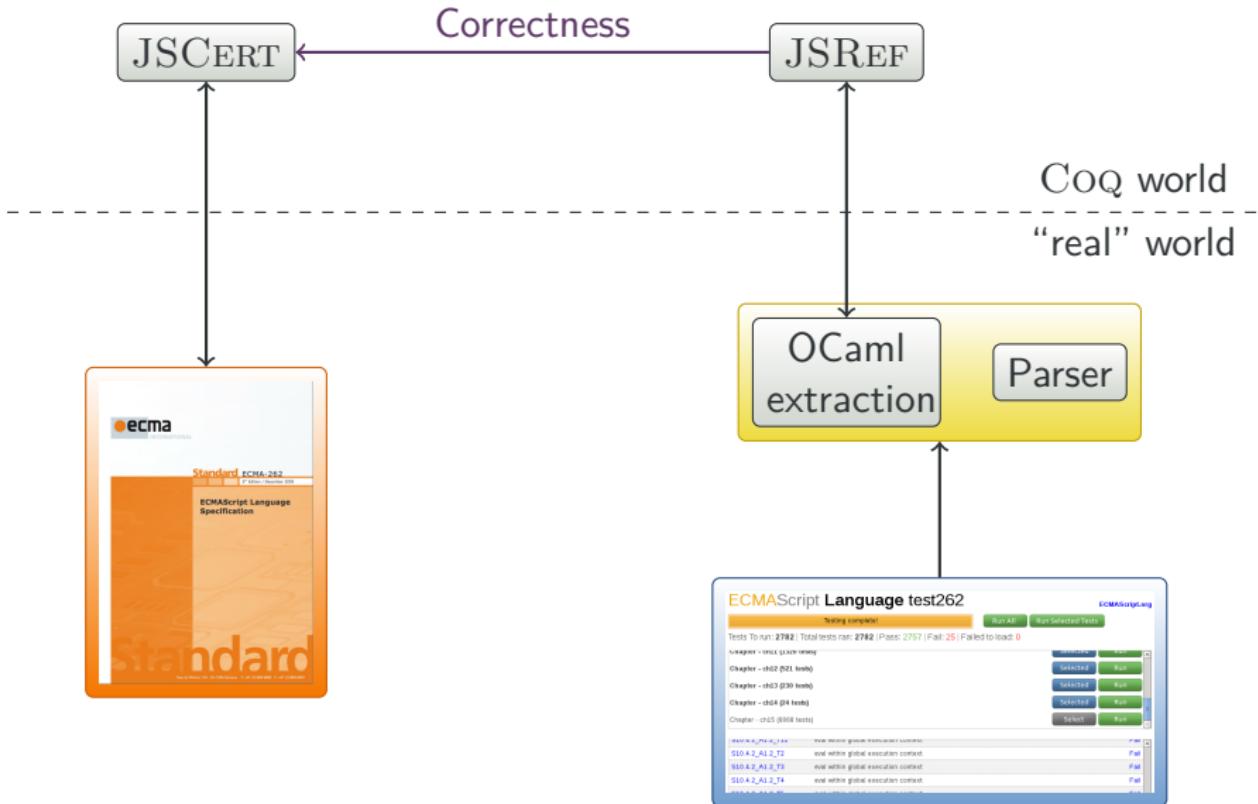
Correctness

If program P **executes** to v, then P and v are **related**

- 2 years, 8 people
- 18 klocs of Coq



Overview



Positive Outcomes

- good coverage of the core of ECMAScript 5.1
- code extraction from JSRef
 - ① instrumented to report coverage
 - ② run the test suite
 - ③ find places not executed (not tested)
 - ④ relate to parts of the spec not tested
 - ⑤ discover discrepancies between implementations

```
1002632 | (** val run_stat_while :  
1002633 |   int -> runs_type -> resvalue -> state -> execution_ctxt -> label_set ->  
1002634 |   expr -> stat -> result **)  
1002635 |  
1002636 | let rec run_stat_while max_step runs0 rv s c ls el t2 =  
1002637 |   (*[77]*)(fun f0 fs n -> (*[77]*)(if n=0 then (*[0]*)(f0 () else (*[77]*)(fs (n-1)  
1002638 |       (fun _ ->  
1002639 |         (*[0]*)(Coq_result_bottom)  
1002640 |       (fun max_step' ->  
1002641 |         (*[77]*)(let run_stat_while' = run_stat_while max_step' runs0_in  
1002642 |           (*[77]*)(if success_value runs0 c (runs0.runs_type_expr s c el) (fun s1 v1 ->  
1002643 |             (*[75]*)(if convert_value_to_boolean v1  
1002644 |               then (*[59]*)(let rR = r2.res_value in  
1002645 |                 (*[59]*)(let rv' =  
1002646 |                   (*[59]*)(let rv' =  
1002647 |                     if resvalue_comparable rvR Coq_resvalue_empty then (*[5]*)(rv else (*[54]*)  
1002648 |                       in  
1002649 |                         (*[59]*)(if_normal_continue_or_break (Coq_result_out (Coq_out_terminator (s2,  
1002650 |                           r2))) (fun r -> (*[41]*)(res_label_in r ls) (fun s3 r3 ->  
1002651 |                             (*[40]*)(run_stat_while' rv' s3 c ls el t2) (fun s3 r3 ->  
1002652 |                               (*[14]*)(Coq_result_out (Coq_out_terminator (s3, (res_ref rv')))))  
1002653 |                             else (*[16]*)(Coq_result_out (Coq_out_terminator (s1, (res_ref rv')))))  
1002654 |                         max_step'))))
```

Scaling Issues

- Hard to keep pace with the standardisation
 - need to update two formalizations and a correctness proof
- JSCert inductive definition is too big
 - no inversion possible, preventing most proofs

Lessons

- many low hanging fruits from an implementation close to the spec
- maintain a single artefact, derive other formats from it
- the Coq formalization should be usable for proofs

To JSExplain

An OCaml interpreter of JavaScript

- very close to the specification
- based on the extraction from JSRef
- uses a tiny subset of OCaml in monadic style
 - functions, tuples, shallow pattern matching, records

1. Let lprim be ? ToPrimitive(lval).
2. Let rprim be ? ToPrimitive(rval).
3. If Type(lprim) is String or Type(rprim) is String, then
 - a. Let lstr be ? ToString(lprim).
 - b. Let rstr be ? ToString(rprim).
 - c. Return the string-concatenation of lstr and rstr.
4. Let lnum be ? ToNumber(lprim).
5. Let rnum be ? ToNumber(rprim).
6. Return the result of applying the addition operation to lnum and rnum.

An OCaml interpreter of JavaScript

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```
and run_binary_op_add s0 c v1 v2 =
  let%prim (s1, w1) = to_primitive_def s0 c v1 in
  let%prim (s2, w2) = to_primitive_def s1 c v2 in
  if  (type_compare (type_of (Coq_value_prim w1)) Coq_type_string)
  || (type_compare (type_of (Coq_value_prim w2)) Coq_type_string)
  then
    let%string (s3, str1) = to_string s2 c (Coq_value_prim w1) in
    let%string (s4, str2) = to_string s3 c (Coq_value_prim w2) in
    res_out (Coq_out_ter (s4, (res_val (Coq_value_prim (Coq_prim_string (strappend str1 str2)))))))
  else
    let%number (s3, n1) = to_number s2 c (Coq_value_prim w1) in
    let%number (s4, n2) = to_number s3 c (Coq_value_prim w2) in
    res_out (Coq_out_ter (s4, (res_val (Coq_value_prim (Coq_prim_number (n1 +. n2)))))))
```

Compiled to JavaScript

- motivations: run it in a browser
- uses compiler-libs to generate a typed AST, which we translate
- target is a tiny subset of JS
 - functions, objects (no prototype), arrays, string, numbers

```
var run_binary_op_add = function (s0, c, v1, v2) {
    return (if_prim(to_primitive_def(s0, c, v1), function(s1, w1) {
        return (if_prim(to_primitive_def(s1, c, v2), function(s2, w2) {
            if ((type_compare(type_of(Coq_value_prim(w1)), Coq_type_string()))
                || type_compare(type_of(Coq_value_prim(w2)), Coq_type_string())))
            {
                return (if_string(to_string(s2, c, Coq_value_prim(w1)), function(s3, str1) {
                    return (if_string(to_string(s3, c, Coq_value_prim(w2)), function(s4, str2) {
                        return (res_out(Coq_out_ter(s4, res_val(
                            Coq_value_prim(Coq_prim_string(strappend(str1, str2)))))); }));})));
            } else { ... }));}));};
```

and to Pseudo JavaScript

- to be readable while staying close to JavaScript
 - hide state and context
 - monadic extension of var
 - pattern matching
 - hide type changes

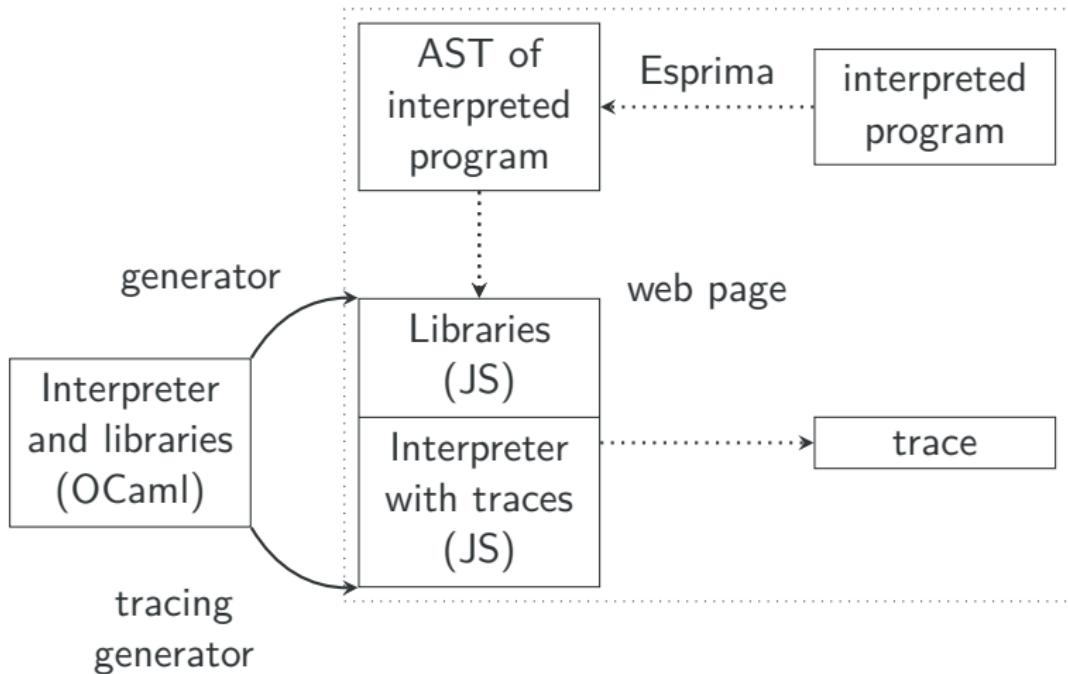
```
var run_expr_binary_op =
function (op, e1, e2) {
  switch (op) {
    case Coq_binary_op_and:
      return (run_binary_op_and(e1, e2));
    case Coq_binary_op_or:
      return (run_binary_op_or(e1, e2));
    default:
      var%run v1 = run_expr_get_value(e1);
      var%run v2 = run_expr_get_value(e2);
      return (run_binary_op(op, v1, v2));
  }
};
```

```
var run_binary_op_add = function (v1, v2) {
  var%prim w1 = to_primitive_def v1;
  var%prim w2 = to_primitive_def v2;
  if ((type_cmp(type_of(w1), Type_string)
    || type_cmp(type_of(w2), Type_string))) {
    var%string str1 = to_string w1;
    var%string str2 = to_string w2;
    return (str_app(str1, str2));
  } else {
    var%number n1 = to_number w1;
    var%number n2 = to_number w2;
    return (n1 + n2);
  }
};
```

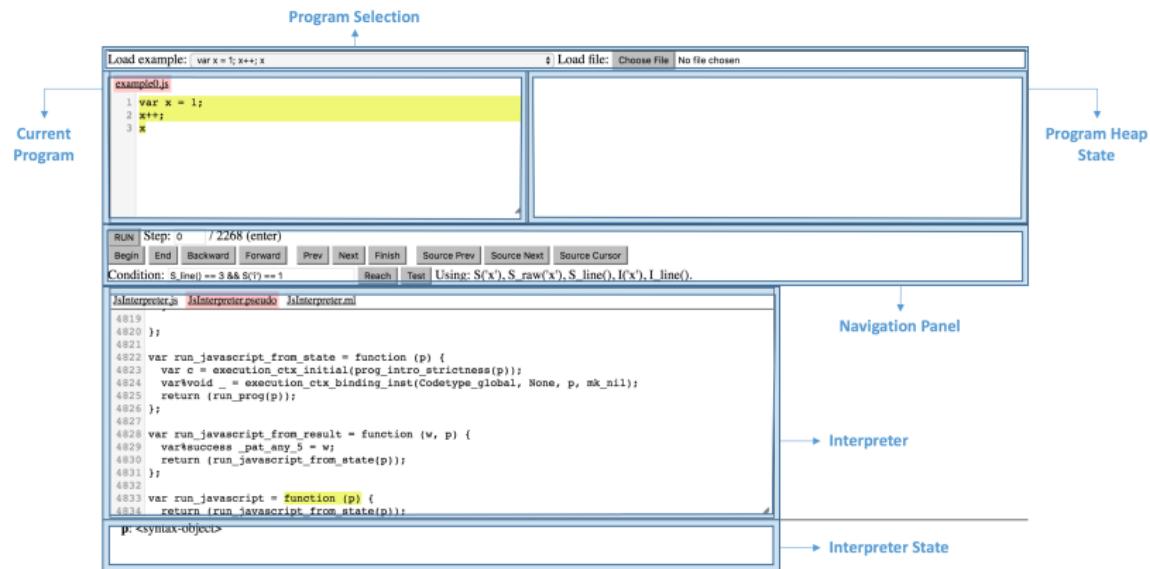
JSExplain

- instrument the generated JavaScript to record *events*
 - Enter (enter a function)
 - CreateCtx(ctx) (new function scope)
 - Add(ident,value) (let binding)
 - Return (return from a function)
- executing the instrumented interpreter generates a trace of events
- web tool to navigate these traces

Architecture



Demo



And Beyond

JSExplain

- extension to current version of JavaScript
 - ongoing, we now can debug it using jsexplain itself
 - engineer hired to work on this in September
- towards a typed specification?
 - PR 1135: Explicitly note mathematical values
 - Issue 496: abstract operations don't always return Completion Records
- better trace navigator
 - links to the specification

Coq Extraction

- needed to prove invariants of the specification
- modular description of the semantics with a simpler induction principle
 - POC for a small language
- we're hiring!¹ (For postdoctoral and PhD positions.)

¹<https://jobs.inria.fr/public/classic/en/offres/2018-00432>

Generalization to Other Languages

- MLExplain²
- plans to do it for Hop.js
- framework to describe semantics

²<https://github.com/Docteur-Lalla/mlexplain/tree/mlexplain>