

Practical Dynamic Symbolic Execution of Standalone JavaScript

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Mission Statement

- Help find bugs in Node.js applications and libraries
- JavaScript is a dynamic language
 - Don't force it into a static type system, invalidates common patterns
 - Static analysis becomes very hard, many sources of precision loss
- Embrace it and go for dynamic approach



- Similar issues as in x86 binary code
 - No types, self-modifying code
- Most successful methods for binaries are dynamic
 - Fuzz testing
 - Dynamic symbolic execution
- No safety proofs, but proofs of vulnerabilities

```

55      pushq   %rbp
48 89 e5      movq   %rsp, %rbp
48 83 ec 20   subq   $32, %rsp
48 8d 3d 77 00 00 00 leaq  119(%rip), %rax
48 8d 45 f8   leaq  -8(%rbp), %rax
48 8d 4d fc   leaq  -4(%rbp), %rax
c7 45 fc 90 00 00 00 movl  $144, %eax
c7 45 f8 e8 03 00 00 movl  $1000, %eax
48 89 4d f0   movq   %rcx, -16(%rbp)
48 89 45 e8   movq   %rax, -24(%rbp)
48 8b 45 e8   movq   -24(%rbp), %rax
8b 10      movl   (%rax), %edx
48 8b 45 f0   movq   -16(%rbp), %rax
89 10      movl   %edx, (%rax)
8b 75 fc   movl  -4(%rbp), %eax
b0 00      movb   $0, %al
e8 21 00 00 00 callq  33
48 8d 3d 3c 00 00 00 leaq  60(%rip), %rax
8b 75 f8   movl  -8(%rbp), %eax
89 45 e4   movl  %eax, -28(%rbp)
b0 00      movb   $0, %al
e8 0d 00 00 00 callq  13
31 d2     xorl   %edx, %edx
89 45 e0   movl  %eax, -32(%rbp)
89 d0     movl  %edx, %eax
48 83 c4 20 addq  $32, %rsp
5d      popq   %rbp
c3      retq
55      pushq   %rbp
48 89 e5      movq   %rsp, %rbp
48 83 ec 20   subq   $32, %rsp
48 8d 3d 77 00 00 00 leaq  119(%rip), %rax
48 8d 45 f8   leaq  -8(%rbp), %rax
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31 d2     xorl   %edx, %edx
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89 d0     movl  %edx, %eax
48 83 c4 20 addq  $32, %rsp
5d      popq   %rbp
c3      retq
ff 25 86 00 00 00      jmpq   *134(%rip)
4c 8d 1d 75 00 00 00 leaq  117(%rip), %rax
41 53     pushq  %r11
ff 25 65 00 00 00      jmpq   *101(%rip)
90      nop
68 00 00 00 00 pushq  $0
e9 e6 ff ff ff jmp    -26 <__stub_...
```

Dynamic Symbolic Execution

- Automatically explore paths
 - Replay tested path with “symbolic” input values
 - Record branching conditions in "path condition"
 - Spawn off new executions from branches
- Constraint solver
 - Decides path feasibility
 - Generates test cases

```
function f(x) {  
  var y = x + 2;  
  if (y > 10) {  
    throw "Error";  
  } else {  
    console.log("Success");  
  }  
}
```

Run 1: $f(0)$:

PC: true
 $x \mapsto X$

Query: $X + 2 > 10$

PC: true

$x \mapsto X$

$y \mapsto X + 2$

Run 2: $f(9)$

PC: $X + 2 \leq 10$

$x \mapsto X$

$y \mapsto X + 2$

High-Level Language Semantics

- Classic DSE focuses on C / x86
 - Straightforward encoding to bitvector SMT
- High-level languages are richer
 - Do more with fewer lines of code
 - Strings, regular expressions

```
function g(x) {  
  y = x.match(/goo+d/);  
  if (y) {  
    throw "Error";  
  } else {  
    console.log("Success");  
  }  
}
```



Node.js Package Manager

Feature	Count	%
Total Packages	415,487	100.00%
Packages with regular expression	145,100	34.92%
Packages with captures	84,972	20.45%
Packages with no source files	33,757	8.10%
Packages with backreferences	15,968	3.84%
Packages with quantified backreferences	503	0.12%

Regular Expressions

- What's the problem?
 - First year undergrad material
 - Supported by SMT solvers: strings + regex in Z3, CVC4
- SMT formulae can include regular language membership

$$(x = \text{"foo"} + s) \wedge (\text{len}(x) < 5) \wedge (x \in \mathcal{L}(/goo+d/))$$

Regular Expressions in Practice

- Regular expressions in most programming languages aren't regular!
- Not supported by solvers

```
x.match(/<([a-z]+)>(.*?)<\/\1>/);
```


Regular Expressions in Practice

- Regular expressions in most programming languages aren't regular!
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`x.match(/<([a-z]+)>(.*?)<\/\1>/);`

capture group

lazy quantifier

backreference

Regular Expressions in Practice

```
x.match(/<([a-z]+)>(.*?)<\/\1>/);
```

- There's more than just testing membership
- Capture group contents are extracted and processed

```

function f(x, maxLen) {
  var s = x.match(/<([a-z]+)>(.*?)</\1>/);
  if (s) {
    if (s[2].length <= 0) {
      console.log("*** Element missing ***");
    } else if (s[2].length > maxLen) {
      console.log("*** Element too long ***");
    } else {
      console.log("*** Success ***");
    }
  } else {
    console.log("*** Malformed XML ***");
  }
}

```

match returns array with matched contents

- [0] Entire matched string
- [1] Capture group 1
- [2] Capture group 2
- [n] Capture group n

- Idea: split expression and use concatenation constraints

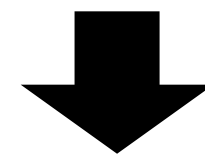
$$t \in \mathcal{L}(/< (a+) > . * ? < \ / \ 1 > /)$$

$$\exists s_1, s_2 : (s_1 \in \mathcal{L}(/a+ /) \wedge s_2 \in \mathcal{L}(/>. *?<\/\ / /) \wedge t = "<" + s_1 + s_2 + s_1 + ">")$$

- Works for membership

- Correct language membership doesn't guarantee correct capture values!

$$t \in \mathcal{L}(/<(a+)>.*?<\/\1>/)$$



$$\exists s_1, s_2 : (s_1 \in \mathcal{L}(/a+/) \wedge s_2 \in \mathcal{L}(/>.*<\/\1>/) \wedge t = "<" + s_1 + s_2 + s_1 + ">")$$

- SAT: $s_1 = "a"; s_2 = "></";$ therefore $t = "<a>"$

Too permissive! Over-approximating matching precedence (greediness)



$\exists s_1, s_2 : (s_1 \in \mathcal{L}(/a+/) \wedge s_2 \in \mathcal{L>(>.*<\/\/)) \wedge t = "<" + s_1 + s_2 + s_1 + ">"$)

- SAT: $s_1 = "a"; s_2 = "></";$ therefore $t = "<a>"$
- Execute `"<a>".match(/<(a+)>.*?<\/\1>/)` and compare
- Conflicting captures: generate blocking clause from concrete result

$$\wedge (s_1 = "a" \rightarrow s_2 = "></")$$

- SAT, model $s_1 = "aa"; s_2 = "></";$ therefore $t = "<a>"$

**Complete refinement scheme with four cases
(positive - negative, match - no match)**



I didn't mention...

- Implicit wildcards
 - Regex is implicitly surrounded with `. * ?`
- Statefulness
 - Affected by flags
- Nesting
 - Capture groups, alternation, updatable backreferences

```
r = /goo+d/g;  
r.test("good"); // true  
r.test("good"); // false  
r.test("good"); // true
```

```
/((a|b)\2)+/
```

ExpoSE

- Dynamic symbolic execution engine (prototype) [SPIN'17]
 - Built in JavaScript (node.js) using Jalangi 2 and Z3
 - SAGE-style generational search (complete path first, then fork all)
- Symbolic semantics
 - Pairs of concrete and symbolic values
 - Symbolic reals (instead of floats), Booleans, strings, regular expressions
 - Implement JavaScript operations on symbolic values

Evaluation



- Effectiveness for test generation
 - Generic library harness exercises exported functions: successfully encountered regex on 1,131 NPM packages
- How much can we increase coverage through full regex support?
 - Gradually enable encoding and refinement, measure increase in coverage

Coverage Increase

Feature	Improved	%	Coverage	Tests/m
No Support	0	0%	+0%	11.46
Regular	528	46.68%	+3.09%	10.14
Encoding	194	17.15%	+2.3%	9.42
Refinement	63	5.57%	+2.18%	8.7
Overall	617	54.55%	+3.39%	

On 1,131 NPM packages where a regex was encountered on a path

Conclusion

- Symbolic execution of code with ECMAScript regex
 - Encode to classic regular expressions and string constraints
 - CEGAR scheme to address matching precedence / greediness
- Robust implementation in ExpoSE
 - Automatic test generation - test oracles currently offloaded to developers
 - Full support for ES5 node.js, including async, eval, regex

<https://github.com/ExpoSEJS>